

**Recall:** Gödel devised a clever trick for coding the symbols in LA as numbers

The key idea is to enumerate all the symbols in a decidable, one-to-one fashion starting from the number 1

e.g. you could enumerate them as:

(	)	^	v	¬	→	∀	∃	+	x	s	0	1	$x_0$	$x_1$	...
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	...

So to every symbol  $s$ , we assign a unique number  $\lceil s \rceil > 0$  and this is done decidable

If we have a string of symbols  $s_0 s_1 s_2 \dots s_n$  then we will code the string as:

$$2^{\lceil s_0 \rceil} 3^{\lceil s_1 \rceil} 5^{\lceil s_2 \rceil} \dots \pi(i)^{\lceil s_i \rceil} \dots \pi(n)^{\lceil s_n \rceil}$$

where  $\pi(i)$  is the  $i^{\text{th}}$  prime and  $\lceil s_0, \dots, s_n \rceil$  is its Gödel number

**Remark:** There are several things to remember:

- $\pi(i)$  is a p.r. function
- Since we don't use 0 as a number for any  $s$ , we know that the number of primes used in its factorization is  $n+1$
- We know and can tell what the exponent is of any prime in this factorization by using the p.r. function  $\text{exp}(n, i)$  — the function which outputs the exponent of  $\pi(i)$  in the prime factorization of  $n$

**Definition:** If we are given a term  $\tau$  then it is just a string of symbols and so it has a Gödel number  $\lceil \tau \rceil$

**Example:** If  $\tau$  is  $SSO$ , then  $\lceil SSO \rceil$  is:

$$2^{\lceil S \rceil} 3^{\lceil S \rceil} 5^{\lceil O \rceil}$$

The exponents which minimize this Gödel number (just for illustrative purposes) is  $\lceil O \rceil = 1$  and  $\lceil S \rceil = 2$ , which means that  $SSO$  (which we commonly know as the number 2) is at least

$$2^2 3^2 5^1 = 180$$

You can imagine that as the strings get larger, this number blows up very fast

**Definition:** Similarly, for any formula  $\varphi$ , we can view it as a string of symbols and find its Gödel number  $\lceil \varphi \rceil$

**Observation:** We can do a number of things algorithmically and in a p.r. fashion

**Example:** Given a number  $n$ , we can determine if  $n$  codes a string of symbols

This entails seeing if the prime factorization of  $n$  involves all the primes from 2 up to some given prime and then writing

$$n = 2^{k_0} 3^{k_1} \dots \pi(m)^{k_m}$$

We then decode  $k_0, \dots, k_m$  as a string of symbols  $s_0, \dots, s_m$  and now we have a decidable process for determining if the string is a term/formula/neither

**Remark:** 10 cannot code a string of symbols since  $10 = 2^1 3^0 5^1$