

## 3-pass protocol, physical version

- The NYTimes wants to have a way of securely receiving anonymous documents.
- It places a lockable box outside its offices and the anonymous source brings the documents, puts them in the box and puts a lock on the box.
- The editors then go outside and put their own lock on the box.
- Later, the source comes back and removes their lock.
- The editors then take their lock off the box and remove the documents.

## 3-pass protocol, digital version

- Bob, who wants to receive something securely (and potentially anonymously), chooses a large prime  $p$  and publishes it.
- Alice encodes her message as a number  $m < p$  and also picks an exponent  $a$ ,  $0 < a < p$  with  $\gcd(a, p - 1) = 1$ . She posts  $m^a \bmod p$ .
- Bob picks his own exponent  $b$ ,  $0 < b < p$  with  $\gcd(b, p - 1) = 1$  and publishes  $m^{ab}$ .
- Alice knows the multiplicative inverse of  $a \bmod (p - 1)$  and so she is able to post  $m^b \bmod p$ .
- Bob then uses the multiplicative inverse of  $b \bmod (p - 1)$  to determine  $m$ .

# Primality testing, Fermat

- The simplest primality test uses Fermat's little theorem.
- Fermat in the contrapositive, says that if  $0 < a < n$  and  $a^{n-1} \not\equiv 1 \pmod n$  then  $n$  is not prime.
- Unfortunately, there are composite numbers  $n$  such that  $a^n \equiv a \pmod n$  for all  $a$ . In fact, there are infinitely many such numbers.

# Primality testing, Miller-Rabin

## Lemma

*If  $x^2 \equiv y^2 \pmod n$  and  $x \not\equiv \pm y \pmod n$  then  $n$  is composite. Moreover,  $\gcd(n, x + y)$  and  $\gcd(n, x - y)$  are non-trivial factors of  $n$ .*

- Here is the Miller-Rabin algorithm for determining primality probabilistically:
- Suppose  $n$  is odd and greater than 9. Write  $n - 1 = 2^k m$  where  $m$  is odd.
- Now pick  $a < n$  randomly and compute a series  $b_i$  for  $i = 0, \dots, k - 1$  as follows:

$$b_0 \equiv a^m \pmod n, b_1 \equiv b_0^2 \pmod n, \dots, b_i \equiv b_{i-1}^2 \pmod n, \dots$$

- We will guess that  $n$  is prime if  $b_0 \equiv \pm 1 \pmod n$  or if  $b_i \equiv -1 \pmod n$  for any  $i$ . Otherwise we will say that  $n$  is composite.

## Primality testing, Miller-Rabin, cont'd

- Claim: If we say  $n$  is composite we will be correct.
- There is at most a 25% chance that if we say  $n$  is prime then we will be wrong.
- If we repeat this test many times for randomly chosen  $a$  and we always get the answer “prime” then with high probability  $n$  is prime.

# Primality testing, Agarwal-Kayal-Saxena

- There is a primality test which is completely deterministic and polynomial in the number of digits of the given number.
- The problem is that its known run-time is order  $n^6$  although 6 is not known to be best possible.
- In practice, if one needs to determine primality for a given number, one uses the probabilistic algorithms to see if you can prove it is not prime and then tries a series of special purpose primality tests which are not efficient but are good enough for "small" numbers, say ones with fewer than 1000 digits.